Configuring random graph models with fixed degree sequences

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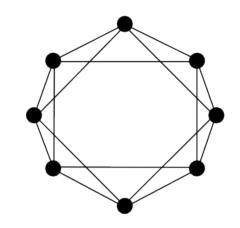


Brief note on references:

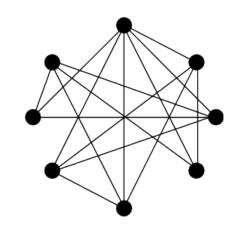
This talk does not include references to literature, which are numerous and important. Most (but not all) references *are* included in the arXiv paper: arxiv.org/abs/1608.00607

Stochastic models, sets, and distributions

a generative model is just a recipe:
 choose parameters→make the network



 a stochastic generative model is also just a recipe: choose parameters→draw a network



- since a single stochastic generative model can generate many networks, the model itself corresponds to a **set of networks**.
- and since the generative model itself is some combination or composition of random variables, a **random graph model** is a set of possible networks, each with an associated probability, i.e., a distribution.

this talk:

configuration models: uniform distributions over networks w/ fixed deg. seq.

Why care about random graphs w/ fixed degree sequence?

Since many networks have broad or peculiar degree sequences, these random graph distributions are commonly used for:

Hypothesis testing:

Can a particular network's properties be explained by the degree sequence alone?

Modeling:

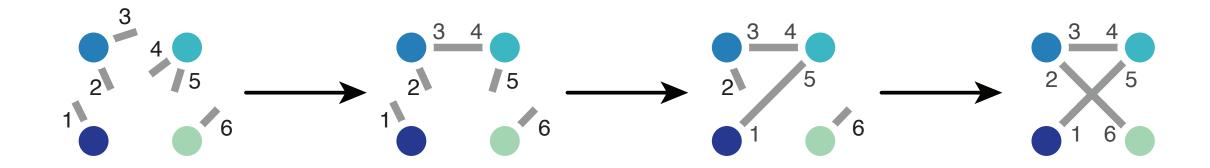
How does the degree distribution affect the epidemic threshold for disease transmission?

Null model for Modularity, Stochastic Block Model:

Compare an empirical graph with (possibly) community structure to the ensemble of random graphs with the same vertex degrees.

Stub Matching to draw from the config. model

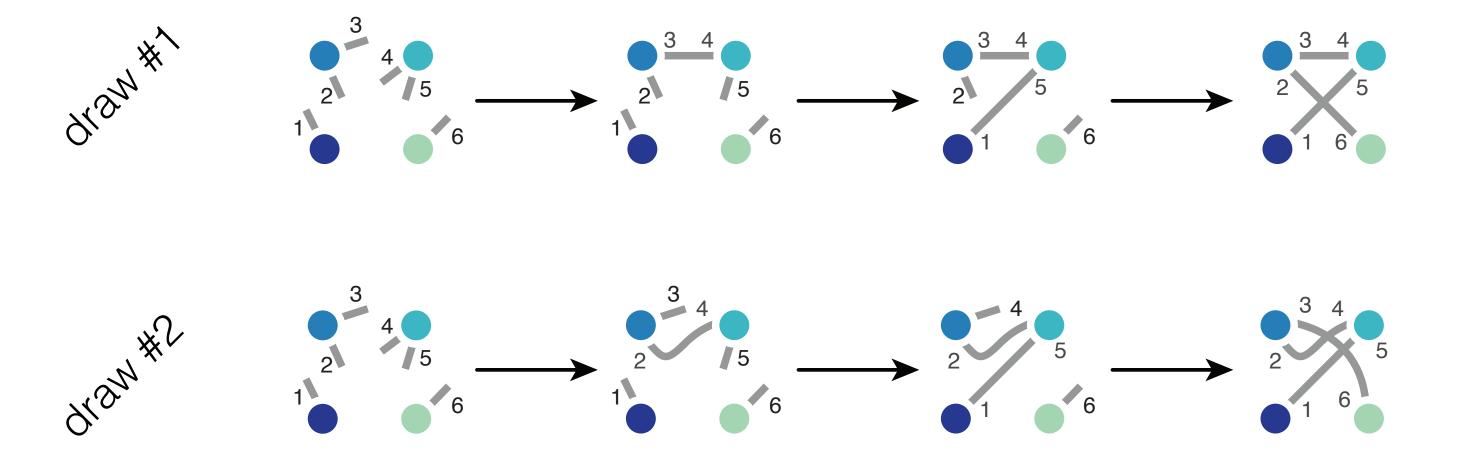
$$\vec{k} = \{1, 2, 2, 1\}$$



the standard algorithm: draw from the distribution by sequential "Stub Matching"

- 1. initialize each node n with k_n half-edges or stubs.
- 2. choose two stubs uniformly at random and join to form an edge.

Stub Matching to draw from the config. model





Are these two different networks? or the same network?

Are stubs distinguishable or not?

The rest of this talk: the answer matters.



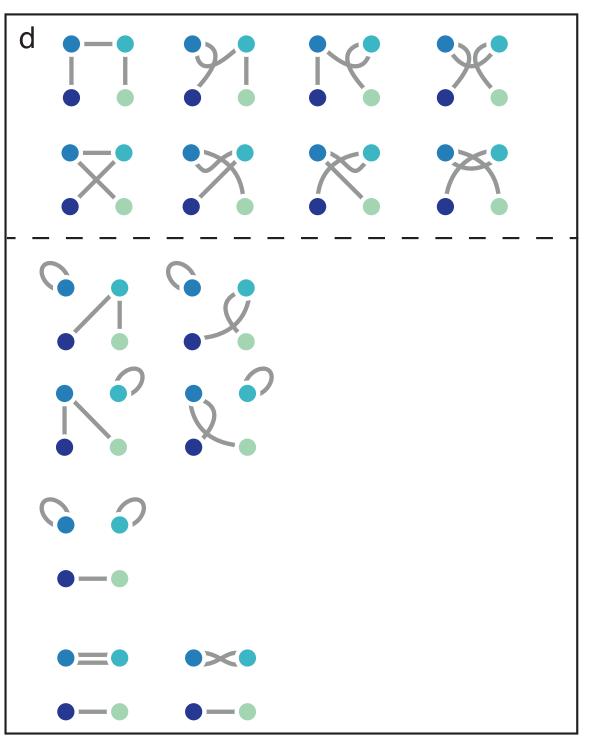
The distribution according to stub-matching

When we draw a graph using stub matching, this is the set of graphs that we uniformly sample.

8 of the graphs are simple, while the other 7 have self-loops or multiedges.

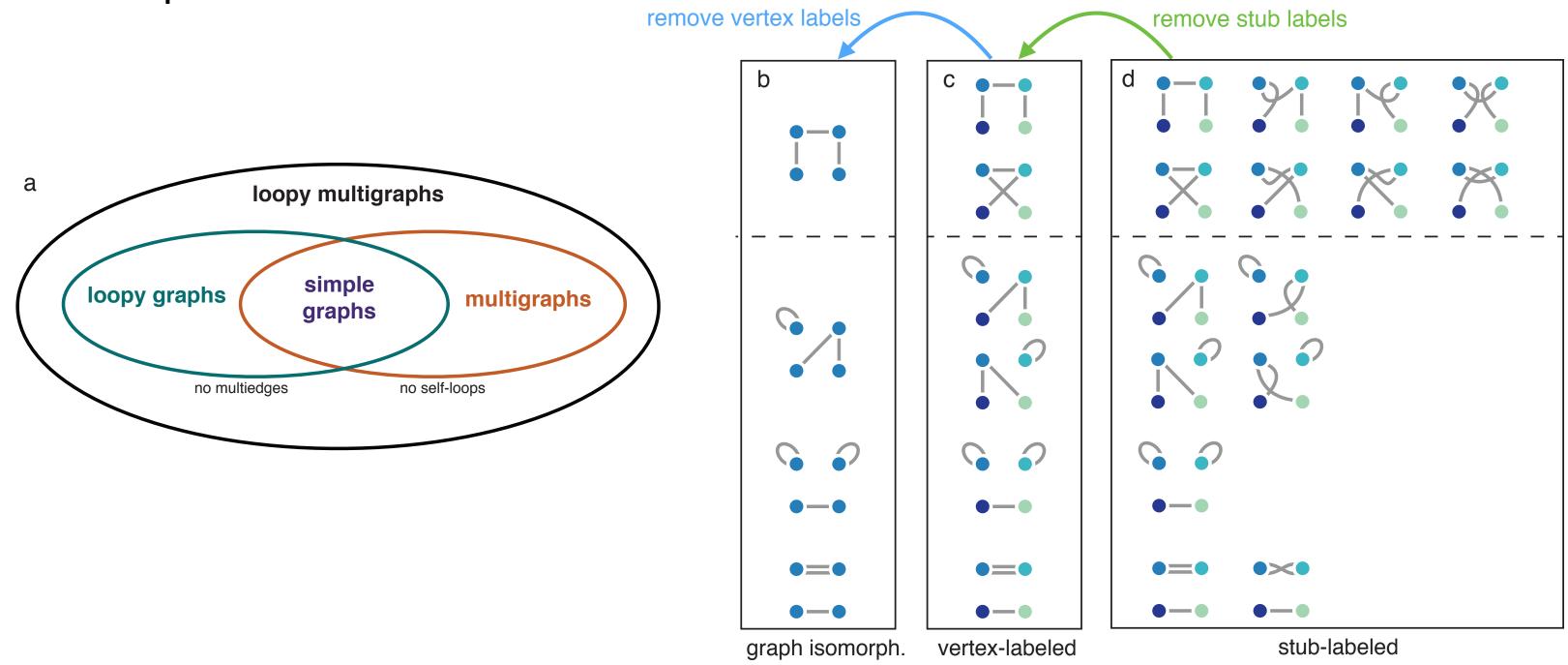
We therefore say that stub matching uniformly samples space of **stub-labeled loopy multigraphs**.

Note, however, that this is not a uniform sample over adjacency matrices (rows).



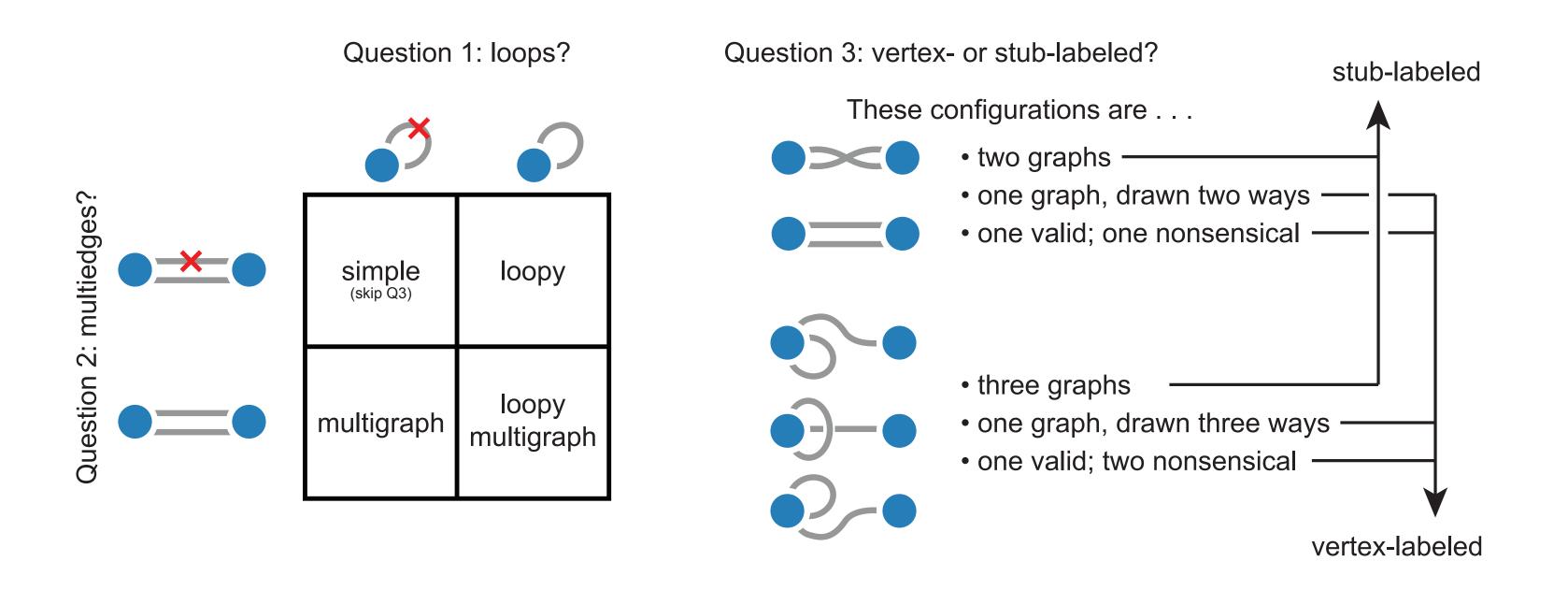
stub-labeled

The importance of uniform distributions



goal: provably uniform sampling for all eight spaces: loopy{0,1} x multigraph{0,1} x {stub-,vertex-}

Choosing a space for your configuration model

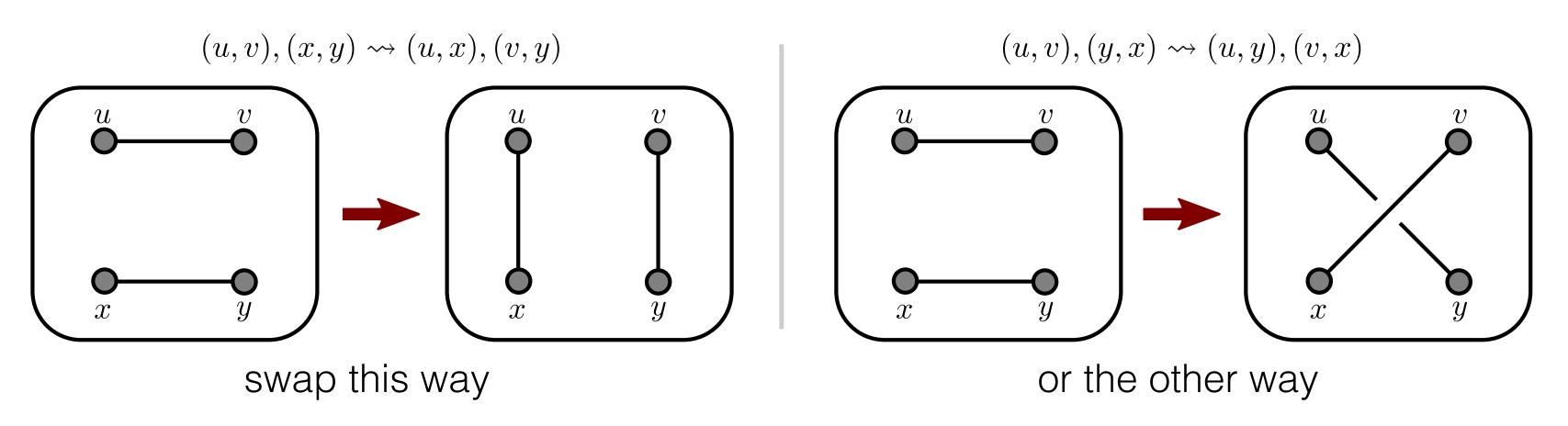


example: Are loops reasonable? Would a loop make sense? [tennis matches: no | author citations: yes]

Sampling from configuration models

stub matching samples uniformly from stub-labeled loopy multigraphs

for other spaces, define a Markov chain over the "graph of graphs" **G**→each vertex is a graph, and directed edges are "double-edge swaps"



NB: Sampling is easy. Provably uniform sampling is not!

Markov chains for uniform sampling

Prove that:

- the transition matrix is doubly stochastic (**G** is regular)
- the chain is irreducible (**G** is strongly connected)
- the chain is aperiodic (G is aperiodic; gcd of all cycles is one)

Straightforward for stub-labeled loopy multigraphs.

Choose two edges uniformly at random and swap them. Accept all swaps and treat each resulting graph as a sample from the U distribution. (Each node in **G** has degree m-choose-2.)

Easy for stub-labeled multigraphs.

Choose two edges uniformly at random and swap them. Reject swaps that create a self-loop and *resample* the current graph. (Think of any "rejected swap" as a self-loop in **G**.)

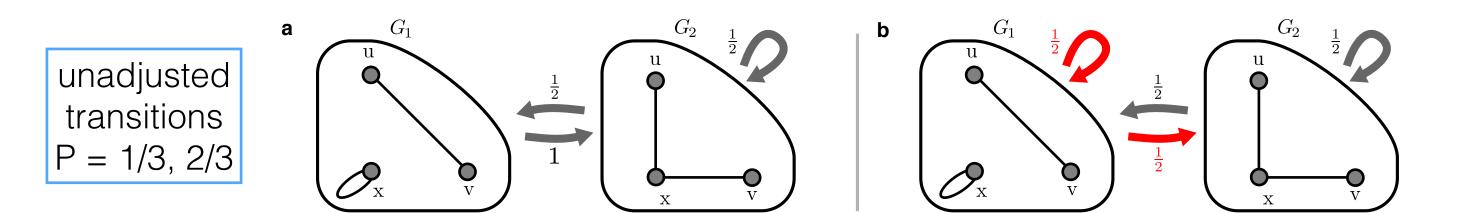
Easy for simple graphs.

Choose two edges uniformly at random and swap them. Reject swaps that create a self-loop or multiedge and *resample* the current graph. (Again, treat "rejected swaps" as a self loops in **G**.)

Markov chains for uniform sampling

For vertex-labeled graphs, we inherit the strong connectedness of **G** as well as its aperiodicity.

However, ensuring that the Markov chain has a uniform distribution as its stationary distribution requires that we adjust transition probabilities.



adjusted transitions P = 1/2,1/2

These asymmetric modifications to transition probabilities depend on the number of self-loops and multiedges in the current state.

Intuition:

decrease outflow (and increase resampling) of graphs with multiedges or self-loops.

Stub-labeled loopy graphs: not connected



counterexample: no double-edge swap connects these two graphs!

Do {stub labels, self-loops, multiedges} matter for how we sample CMs? yes

showed that these spaces are far from equivalent, even in thermodynamic lim.



introduced (and just outlined) provably uniform sampling methods.

Do (stub labels, self-loops, multiedges) matter in applications of CMs? next...

→hypothesis testing

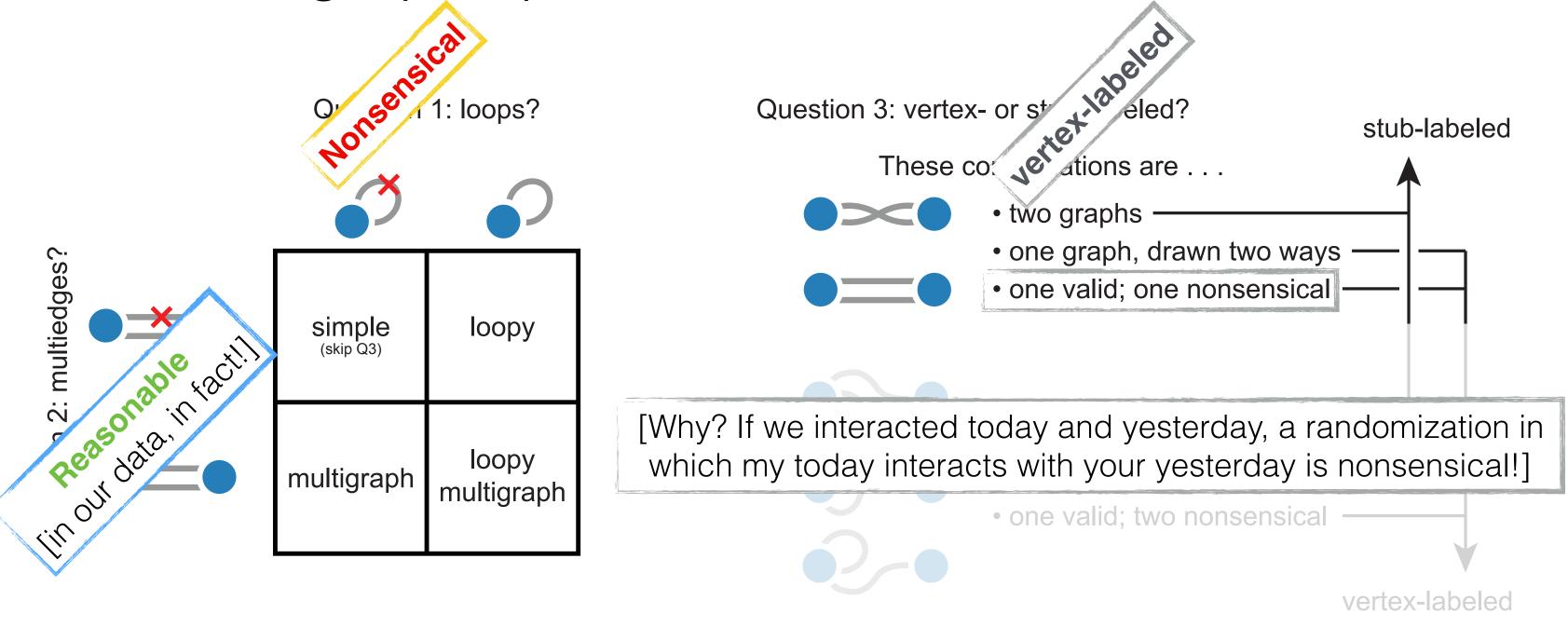


→null model for modularity

Hypothesis testing

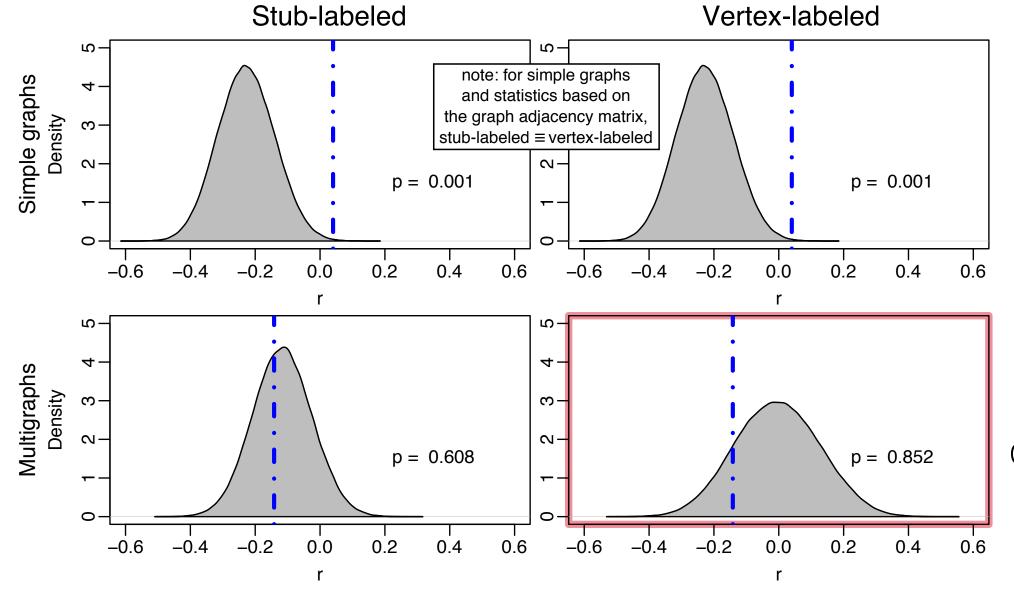


Choose a graph space for barn swallows



This should be modeled as a vertex-labeled multigraph.

Assortative pairing of barn swallows



Sanity check: should be = for simple

NONE of these is centered at zero. Correct space is *meaningfully* different.

Uniform sampling means we can compare empirical value to null distribution to draw scientific conclusions.

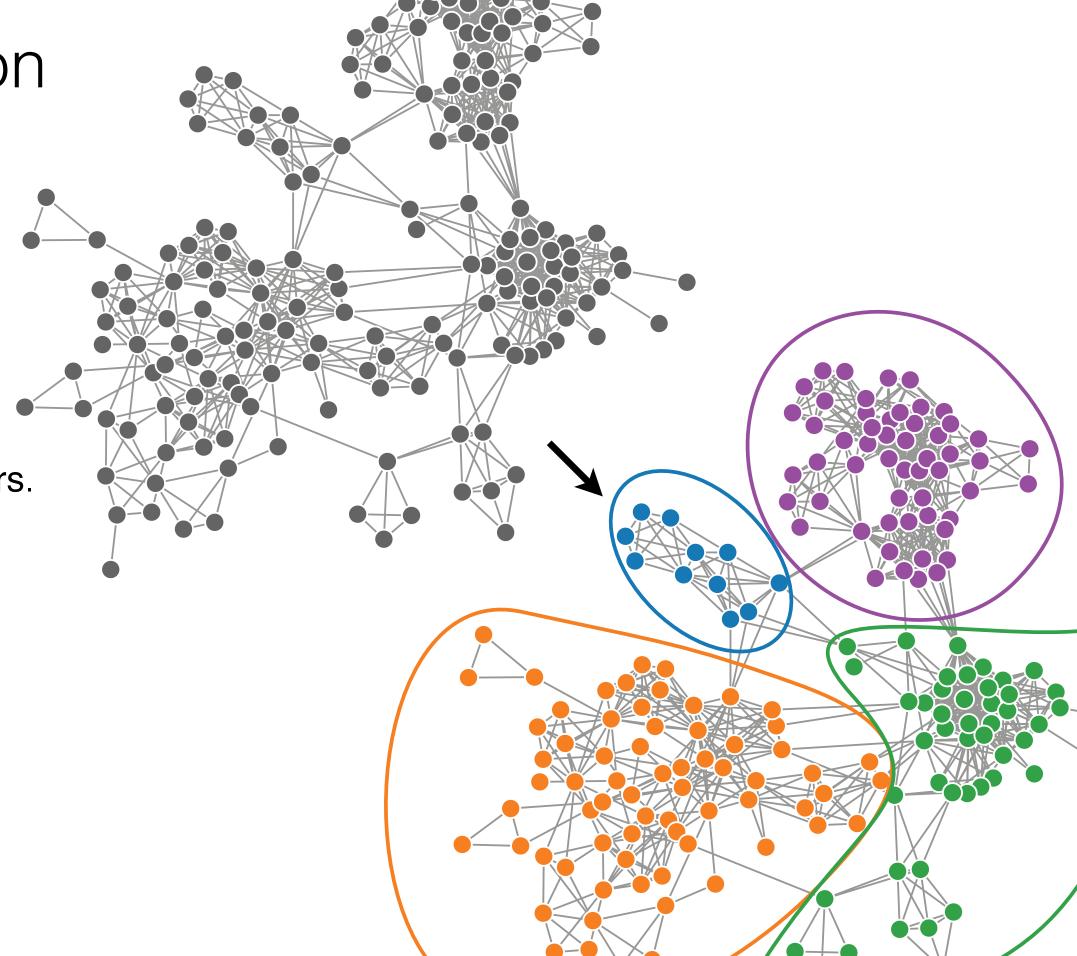
The choice of graph space matters—careful choice & sampling can flip conclusions!

Community Detection

Are there groups of vertices that **tend to associate** with each other more than **we expect by chance**?

Data: collaborations among geometers.

Maximize modularity, e.g.



Coauthorship communities (vertex-labeled multigraph)

random

expected number of edges in a random degree-preserving null model

Modularity

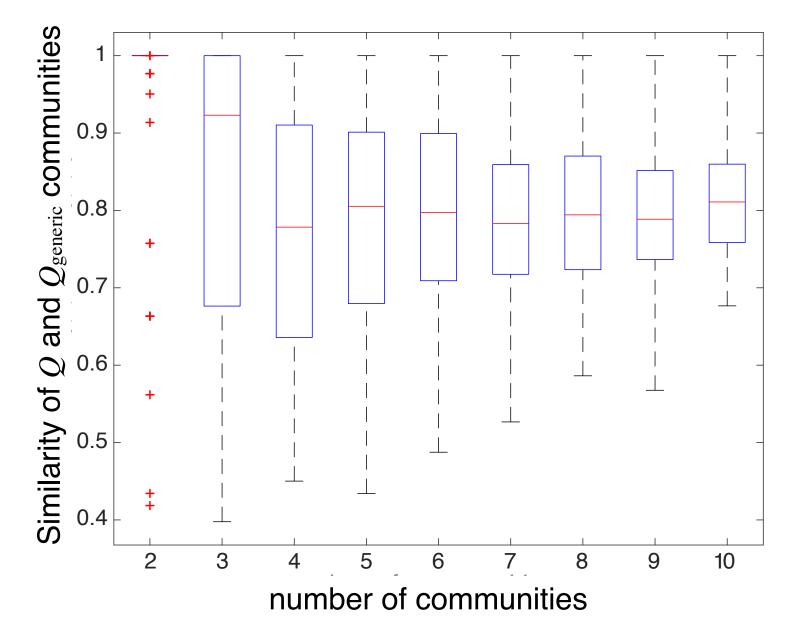
$$Q = \frac{1}{2M} \sum_{i,j} \left(A_{ij} - \frac{k_i k_j}{2M} \right) \delta(g_i, g_j)$$

specifically, in the stub-labeled loopy multigraph CM

Generic Modularity

$$Q_{\text{generic}} = \frac{1}{2M} \sum_{i,j} \left(A_{ij} - \frac{\mathbb{E}[C_{k_i,k_j}]}{n_{k_i} n_{k_j}} \right) \delta(g_i, g_j)$$

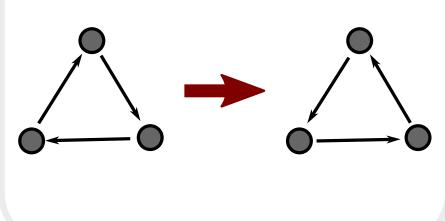
expected number of edges in *any* random degree-preserving null model



same community detection algorithm, same initial state, different results

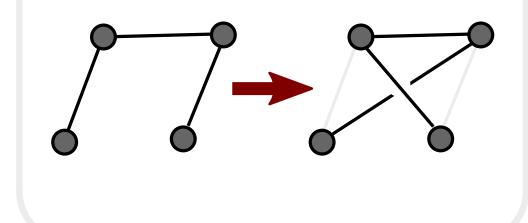
Advanced edge swaps

a reversing a directed triangle

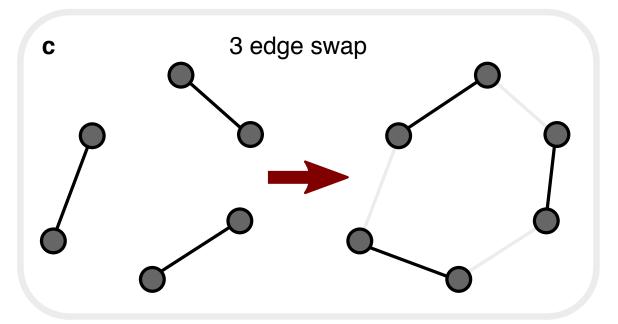


required for graph-of-graphs irreducibility in directed networks

b connectivity preserving edge swap



useful if you wish to sample *only* networks that have a fixed number of connected components



other swaps have been proposed, e.g. to improve mixing time

Proofs, samplers, the history of the configuration model, and applications in the paper

The point: graph spaces & stub labels *matter*, in theory and in practice.

Recognizing this exposes a number of unrecognized & unsolved problems.

Provably uniform sampling methods exist—some have existed for decades!







Joel Nishimura





Configuring Random Graph Models with Fixed Degree Sequences Fosdick, Larremore, Nishimura, Ugander. To Appear in *SIAM Review.* arxiv.org/abs/1608.00607



danlarremore.com/configuration models





Thank you

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